## HP Forums / HP Calculators (and very old HP Computers) / General Forum $\nabla$ / [VA] SRC \#012b - Then and Now:

 Root
## Valentin Albillo 8

Senior Member

Posts: 958
Joined: Feb 2015
Warning Level: 0\%

## [VA] SRC \#012b - Then and Now: Root

## Hi, all,

After the many excellent solutions \& comments posted for Problem 1 and once the $\mathbf{7 , 1 0 0}$ views mark has been met and exceeded, now's the time for the $2^{\text {nd }}$ part of my new SRC \#012 - Then an Now, where I'll demonstrate that advanced vintage HP calcs which were great problem-solvers back THEN in the 80's (some 40 years ago !) are NOW still perfectly capable of solving recently-proposed tricky problems intended to be tackled by using fast modern 2020-era personal computers, never mind slow ancient pocket calcs.

In the following weeks I'm proposing six increasingly harder such problems for you to try and solve while respectfully abiding by the following mandatory rules summarized here:

> You must try and solve them using EXCLUSIVELY your preferred VINTAGE HP CALCULATORS (physical or virtual,) coding in either RPN (including HP-41 microcode), RPL (any variant existing at the time, including SysRPL) or HP-71B languages (including BASIC, FORTH and Assembler) AND NOTHING ELSE: NO CAS/XCAS, MATLAB, MATHEMATICA, EXCEL, C/C++/C\#, PYTHON or the like, AND NO LENGTHY MATH SESSIONS/LECTURES. Also, NO CODE PANELS, please!
> On the positive side, you may use any official/well-known modules, pacs or libraries which were available at the time, such as the Math Pac and JPC ROM for the HP-71B, the Advantage Module, PPC ROM and Extended Memory for the HP-41, and assorted libraries for the RPL models, to name a few.

This Problem 2 deals with polynomial roots with a bang, namely:

## Problem 2: Root

Write a program to find the minimum absolute value among the roots of the following polynomial:

$$
P(x)=2+3 x+5 x^{2}+7 x^{3}+11 x^{4}+13 x^{5}+\ldots+104743 x^{10,000}
$$

whose coefficients are the prime numbers in order: $2,3,5,7,11,13, \ldots, 104743$.

Your program should have no inputs and must output the asked value and automatically end. You should strive for 10-12 correct digits (gave or take a few ulp) depending on your HP model, and the faster the running time the better. Also, you must justify in your comments the soundness of your approach, not "just trying" or relying on luck.

Some useful advice is to try and find the correct balance between letting the program do all the work (i.e. sheer brute force, which could potentially take far too much RAM and running time) with no help from you, or else use a little bit of insight to help significantly speed up the process. Your choice.
(As an aside, I wonder if any (or even all !) forum members who successfully posted correct solutions for Problem 1 will be able to follow suit and solve this Problem 2 as well, for a perfect 2 for 2 score! (ص))

If I see interest I'll post in a few days my original solution for the HP-71B, a 5-liner which computes the required absolute value relatively quickly and accurately (it can be done in just 4 lines albeit at a significantly slower speed).

In the meantime, let's see your very own clever solutions AND remember the above rules, please.
$v$.

Posts: 767
Joined: Dec 2013

## RE: [VA] SRC \#012b - Then and Now: Root

(second try..)
Well I do have a result now, my previous one was wrong.
It's on Free42, which appears to be permitted (I understood that wrongly), and while I can make it run on a 42 S , it would take too long.
Also, I have no idea if my result is the smallest. I'll have to do some thinking instead of mindless coding ;-)
Werner

- the Free42 solution precalculates the matrix of 10001 coefficients - an easy adaptation of a prime number listing routine I wrote years ago. On a 42S, this is not feasible, so the coefficients would have to be calculated over and over - hence the probably long execution time there. (it takes about 4 hours to generate the coefficients once)
- then, a routine to evaluate a polynomial stored as a row vector, also from my archives (really simple)
- and at last, a 'specific' solver routine, which, I think, was inspired by a post of Valentin himself, long ago. 7 iterations are needed to converge to a result, so, no, I'm not going to run it on my 42S.
When I'm reasonably sure the result is correct I'll post some code.


## J-F Garnier 8 <br> Senior Member

Posts: 790
Joined: Dec 2013

## RE: [VA] SRC \#012b - Then and Now: Root

I was not sure to fully understand the problem. Do you have to search for the smallest (in absolute value) root in the real or complex domain?
When we are speaking of roots of a polynomial, we are often referring to its complex roots.
But when I read 'absolute value', I tend to link it to a real number, even if the ABS function is often used in HP programming languages to calculate the norm or modulus of a complex number.

So I will assume that we are looking for real roots.
My first thoughts: a real root must be negative, and greater than -1 since the polynomial quickly takes very large values for $|X|>1$.

My first analysis (using HP calculators), cutting the polynomial to the first 10-100 terms seemed to indicate me that there is no real root.
So if there are real roots, they are close to -1 , otherwise the higher $\mathrm{X}^{\wedge} \mathrm{i}$ terms would be negligible and I would have found the root with the truncated polynomial.
I will try to build a program to support this, but I'm not sure I will be able to propose something useful, even on a fast emulator.

J-F

```
5コ:5コ:5` Werner 8 Posts:767
Senior Member Joined: Dec 2013
```


## RE: [VA] SRC \#012b - Then and Now: Root

(error corrected in PX)
There are indeed no real roots.
I simply built the polynomial, stored its coefficients in a matrix $P$, and used my complex solver (explanation and source code here) with starting value 0 to find
( $-0.645758096347,0.483177676217$ ), abs value 0.806513599261
here's the polynomial generator routine:

```
00 { 139-Byte Prgm }
01•LBL "PLST"
02 1
0310001
04 NEWMAT
05 EDIT
06 1.001
07 STO 00
0 8 \text { SIGN}
```

09 STO IND 00
102
11 XEQ 14
123
13 XEQ 14
$14+$
15 XEQ 14
$16 \cdot \mathrm{LBL} 02$
172
18 XEQ 03
19 FS? 77
20 GTO 00
214
22 XEQ 03
23 FC? 77
24 GTO 02
$25 \cdot \mathrm{LBL} 00$
26 RCLEL
27 EXITALL
28 STO "P"
29 RTN
$30 \cdot$ LBL 03
$31+$
32 RCL 00
33 X<>Y
$34 \cdot L B L 04$ @ ( loop over all stored primes )
35 RCL IND ST Y
36 RCL ST Y
$37 \cdot L B L 05$ @ (GCD )
38 MOD
39 LASTX
40 X $<>Y$
$41 X>0$ ?
42 GTO 05
$43+$
44 R $\downarrow$
45 DSE ST T @ ( test whether GCD=1 )
46 RTN
47 ISG ST Y
48 GTO 04
@ ( either the number is prime or it is a perfect square of one )
49 ENTER
50 SQRT
51 IP
$52 X^{\wedge} 2$
$53 X<>Y$
$54 \mathrm{X}>\mathrm{Y}$ ? @ ( number is prime )
55 GTO 14
56 SQRT @ ( number is a perfect square )
57 STO IND ST Z @ ( store it in next register, just in case )
58 DSE ST Z @ ( will always skip )
$59 \mathrm{X}>0$ ? @ (nop)
60 RCL IND ST Z
$61 \times$
62 SIGN @ ( see if product is too large )
63 LASTX
64 STO + ST Y
$65 \mathrm{X}=\mathrm{Y}$ ?
66 GTO 00
67 STO IND ST T @ (if not, save it )
68 RCL ST Z
69 RTN
70•LBL 00 @ ( else increase array )
71 1е-3
$72 \mathrm{STO}+00$
73 R^
74 RTN
75•LBL 14 @ ( separate label to easily change )
$76 \rightarrow$
77 END
and the polynomial evaluator (coefficients [a0 a1 .. an] so Horner scheme has to start at the end)

```
00 { 25-Byte Prgm }
01•LBL "PX" @ evaluate polynomial, matrix indexed at (1,1)
0 2 ~ E N T E R
03 ENTER
04 ENTER
05 CLX
0 6 ~ J - ~ @ ~ s t a r t ~ f r o m ~ t h e ~ e n d
07•LBL 02
0 x
09 RCLEL
10 +
1 1 ~ J -
2 FC? }7
1 3 \text { GTO 02}
1 4 ~ J + ~ @ ~ b a c k ~ t o ~ ( 1 , 1 )
15 END
```

There is, however, a much better way, one that singles out the smallest root (thanks Albert!). Working on that ;-) Cheers, Werner
C.Ret 8
Posts: 223
Member
Joined: Dec 2013

RE: [VA] SRC \#012b - Then and Now: Root
Ah! Ah!

It is with great pleasure and impatience that I discover this new stage of the challenge!
I rub shoulders with it, but I'm worried. The parity of the polynomial is a problem for me as well as the slowness of my HP-71B. Couldn't things be more complicated than they seem?

As I'm not very proud of my performance and the capabilities of my HP-71B (unless it's the other way around?) and in order not to waste too much time developing my solution, I decided to tackle a much simpler polynomial:
$P_{9}(x)=2+3 x+5 x^{2}+7 x^{3}+11 x^{4}+13 x^{5}+17 x^{6}+19 x^{7}+23 x^{8}+29 x^{9}$

This polynomial has same resemblances with the polynomial $\boldsymbol{P}$ which interests us but has the enormous advantage of being infinitely shorter.

Please note that the derivate of $P_{9}$ is easily determinate as:
$P_{9}^{\prime}(x)=3+2 \cdot 5 x+3 \cdot 7 x^{2}+4 \cdot 11 x^{3}+5 \cdot 13 x^{4}+6 \cdot 17 x^{5}+7 \cdot 19 x^{6}+8 \cdot 23 x^{7}+9 \cdot 29 x^{8}$
This allows me to code faster; Using Horner's scheme, Newton's root finding method and the prime number facility built in the JPC ROM cartridge, I manage to compose a code adapted to my HP-71B which in a few tens of seconds gives me the absolute value of the only real solution:

```
10 DESTROY ALL !! SRC#012b V.A. challenge (v.9)
    @ REAL D,E,K,P,X,Y
    @ X=-2/3 @ E=1.E-12 ! Initial guess and accuracy
20 REPEAT
    @ P=29 @ Y=P @ D=0
    !! Computation of Y=P(X) and D=P'(X)
    P last prime 
            D=D*X+K*P
            P=FPRIM(P-1,2) ! FPRIM(P-1,2) return previous prime
            Y=Y*X+P
            NEXT K
40 X=X-Y/D ! Newton's method X(n+1) = X(n) - P(x)/P'(x)
    UNTIL ABS(Y/D)<E
    DISP ABS (X) @ END
                                    !! Display absolute value of
root X
```

$\left|x_{9}\right|=.79462$

To increase speed and easiness, values of $P_{9}$ and its derivate $P_{9}^{\prime}$ are compute together in the same FOR TO NEXT loop.

My next step was to use this exact algorithm for the next polynomial:
$P_{10}(x)=2+3 x+5 x^{2}+7 x^{3}+11 x^{4}+13 x^{5}+17 x^{6}+19 x^{7}+23 x^{8}+29 x^{9}+31 x^{10}$

But I suddenly discover what already discover Werner and J.-F. Garnier; polynomials with only positive coefficients and even degree have no real root. So my previous version of the code didn't converge at all!

No stress, thanks to his Math module, this HP-71B is much more powerful that any basic calculator. It is very easy to modify the program to run the same algorithm but with complex values:

```
1 0 \text { DESTROY ALL !! SRC\#012b V.A. challenge (v.10)}
    REAL E,K,P
    COMPLEX D,X,Y
    X=(.3,.7) @ E=1.E-12 ! Initial guess and accuracy
    REPEAT !! Computation of Y=P(X) and D=P'(X)
    P=31 @ Y=P @ D=0 ! P last prime
    FOR K=10 TO 1 STEP -1 Y/D computation loop
                D=D* X +K* P
                P=FPRIM(P-1,2) ! FPRIM(P-1,2) return previous prime
                Y=Y*X+P
        NEXT K
        X=X-Y/D ! Newton's method X (n+1) = X (n) - P(x)/P'(x)
    UNTIL ABS (Y/D)<E
    DISP ABS(X) @ END !! Display absolute value of
root X
```

For $P_{10}$, I easily get the following absolute minimal value (in fact norm of the complex root).
$\left|z_{10}\right|=.734576$ since the closest complex root I found for $P_{10}$ near $(0,0)$ is $z_{10}=(0.297370,0.671694)$
P.S.: Using this complex valued algorithm with $P_{9}$ clearly indicate that I have not found the minimum absolute value among the roots since other complex roots exists that potentially have a smaller absolute norm.
For instance: $\left|z_{9}\right|=0.71297$ due to root $(-.56933, .42917)$ or conjugate.

Now, I need new batteries and a large bunch of time to run my code for the true challenge for the lengthy polynomial! 10 '000 that's really huge!
I urgently need to find a way to determine the correct initial guess at first attempt.
References: Here is one of the documents that inspire me and help me develop this exact solution.
Newton's method, and the fractal it creates that Newton knew nothing about. At 11:16 start the chapter about the 'Fun Facts', as fun as this Valentin Challenge. But you may watch there why I am stuck at the moment and need another good idea to efficiently start a 3 -hour computation on the right initial guess!


Posts: 790
Joined: Dec 2013

RE: [VA] SRC \#012b - Then and Now: Root
Thanks Werner and C.Ret to put me back on the right track, so we are looking for complex roots.
But I didn't completely loose my time, just see:

## Werner Wrote:

(9th November, 2022 11:02)
There are indeed no real roots.

## C.Ret Wrote:

(9th November, 2022 12:19)
But I suddenly discover what already discover Werner and J.-F. Garnier; polynomials with only positive coefficients and even degree have no real root.

This is wrong! There are (at least) two real roots !

I used brute force, using HTBasic (1999 version) which is a HP BASIC compatible programming environment on PC, that still runs fine on my W10 computer in 2022.
I know this is slightly outside the rules, but not so much, HTBasic is to the HP-71B what Free42 is to the HP-42S: a native (not emulated) compatible language running on PC, with the notable difference is that HTBasic is not free, actually quite expensive being a professional tool.
So the below program, that searches where the polynomial sign changes, could in principle, be run on a HP-71B:

```
100 ! RE-STORE "src12b"
05 ! SRC12B, MoHPC, 8nov2022
110 ! quite 'force brut' approach
115 !
120 OPTION BASE 0
DIM P(10000)
!
N=10000
READ P(*) ! READ P() on the 71B
!
! find out where the sign changes:
Y1=2 ! polynomial value at X=0
FOR X=-.01 TO -1.01 STEP -.001
    GOSUB Evalpoly
    IF SGN(Y) <>SGN(Y1) THEN
            PRINT X1,Y1 ! print the interval where sign changes
            PRINT X,Y
            PRINT "------"
        END IF
        X1=X ! save X,Y for next iteration
        Y1=Y
    NEXT X
    STOP
!
Evalpoly: ! evaluate polynomial at X, result in Y
Y=0
FOR I=0 TO N
    IF (I*LGT(-X))>-300 THEN Y=Y+X^I*P(I)
NEXT I
RETURN
!
290!
295 DATA 2, 3, 5, 7, 11, 13, 17, 19, 23, 29
300 DATA 31, 37, 41, 43, 47, 53, 59, 61, 67 , 71
305 DATA 73, 79, 83, 89, 97, 101, 103, 107, 109, 113
5290 DATA 104677 , 104681 , 104683, 104693 , 104701, 104707 , 104711 , 104717 , 104723 , 104729
5 2 9 5 \text { DATA 104743}
5300 END
```

X, polynomial value:
-. 996 . $9768 .$.
-. 997 -5.9593...
------
$-.999-39.8287 \ldots$
-1 52726
-

So there is a real root between -. 996 and -.997 and another one between -.999 and -1 , in accordance to my guess that real roots (if existing) would be close to -1 .

Now, I will go back to the actual question, using the 71B or a HP calculator (promised!), but I thought this unexpected result was worth to be reported here.

J-F
C.Ret 8

Member

X, polynomial value:
-. 996 . 9768 ...
-. 997 -5.9593...
-------
-. 999 -39.8287...
-1 52726
-----
So there is a real root between -.996 and -.997 and another one between -.999 and -1 , in accordance to my guess that real roots (if existing) would be close to -1 .

That's of great interest. I am compiling a short HP-71B -Math Pack JPC ROM HP-BASIC code to confirm your observations. This will also indicate me how slow is my device to compute these only 4 values.

But that won't allay my worries; among the 9998 other complex roots, is there not one or the other closer to the point $(0,0)$ ?

* QUOTE की REPORT

Posts: 172
Joined: Jul 2015
C.Ret Wrote:
(9th November, 2022 18:37)
So there is a real root between -.996 and -.997 and another one between -.999 and -1 , in accordance to my guess that real roots (if existing) would be close to -1 .

That is very perplexing. At first I immediately gave up given then large number of coefficients and powers thinking of the solve algorithm we have for the hp41. But then I decided to just play around a little bit, using simpler polynomials.

I made my way in a somewhat (ok, very) manual sleuthing fashion across the first 30 or so polynomials. And then looked at the real roots. In all cases, I was not able to find a real root for a polynomial which ended in a even power. Which peter-intuitively (ie inferior, wrong, intuition) made sense - the last term has the highest exponent and highest multiplier and as such dominates the prior term. And this is true for all pairs of "even, odd" powers. So each pair creates a positive overhang, making it impossible to converge. Or so my logic went. Clearly, contra factum non est discudandum so I need to think more about this (or get some help from the team).

Using the real roots for the odd-ending polynomials, I saw a really nice curve that looked like a logarithmic curve. And low and behold, a logarithmic fit gives some $92 \% R^{\wedge} 2$. And would point to a solution slightly bigger than -1 . So I need to think about why that pretty smooth and steady fit for the odd-ending polynomials would have a minimum and turn back up. Only thing I can think of is the space between primes is getting larger, creating bigger gaps between the even-odd pairs.

However, all my hopes sank when I realized that we are looking for the smallest absolute solution. And there is no way I could think of (so far) to find the smallest absolute value directly, out of 5000 pairs of conjugated roots for the full polynomial.

I will think how to do that in the smaller polynomial case.
it is also possible that I can find a way to show that the 10 digit precision of the hp41 only goes to a much smaller polynomial and as such I can just solve a much smaller polynomial as an approximation inside the accuracy of the HP chosen.


## J-F Garnier 8

Senior Member

Posts: 790
Joined: Dec 2013

## RE: [VA] SRC \#012b - Then and Now: Root

## PeterP Wrote:

(9th November, 2022 22:13)
I made my way in a somewhat (ok, very) manual sleuthing fashion across the first 30 or so polynomials. And then looked at the real roots. In all cases, I was not able to find a real root for a polynomial which ended in a even power. Which peter-intuitively (ie inferior, wrong, intuition) made sense - the last term has the highest exponent and highest multiplier and as such dominates the prior term. And this is true for all pairs of "even, odd" powers. So each pair creates
a positive overhang, making it impossible to converge. Or so my logic went. Clearly, contra factum non est discudandum so I need to think more about this (or get some help from the team).

Let's take the last two terms:
104729 * X^9999 + 104743 * X^10000
for $X=-1$ you get +14
but for $X=-0.999$ you get -0.0041
do the sum of 5000 such pairs, and you can get something quite negative, large enough to overcome the constant 2 factor.
Matter of fact, I found real roots for polynomials of 5000 and 10000 terms, but not with only 2000 terms or less.
J-F

## Fernando del Rey 8

Posts: 19
Junior Member
Joined: Dec 2013

## RE: [VA] SRC \#012b - Then and Now: Root

I've written a short program for the HP-71B that will solve the problem for polynomials of degree up to a few hundred, using brute force with the Math ROM's PROOT function.

What I found out is that, for polynomials of degree 149 and beyond, the roots with the minimum absolute value have always an absolute value of:

### 0.80651359926

I have tried with polynomials of degree up to several hundred always getting the same minimum absolute value, so I am assuming that this value would be the same for the polynomial of degree 10000 in Valentin's OP. But I don't have a hard proof for it, it's just a guess.

I ran my program using Emu71/Win, where it takes just 3 or 4 seconds to get the result for degree 150 . On a real 71 B it would take several minutes, but I haven't measured it.

## J-F Garnier 8

Senior Member

Posts: 790
Joined: Dec 2013

RE: [VA] SRC \#012b - Then and Now: Root

## Fernando del Rey Wrote:

(10th November, 2022 01:03)
I've written a short program for the HP-71B that will solve the problem for polynomials of degree up to a few hundred, using brute force with the Math ROM's PROOT function.

What I found out is that, for polynomials of degree 149 and beyond, the roots with the minimum absolute value have always an absolute value of:

### 0.80651359926

I did the same yesterday evening with 200 terms, and found the same minimum root as you and Werner already found: $Z \mathrm{~min}=(-0.645758096347,0.483177676217)$ for an abs value of 0.806513599261

My understanding is that the PROOT method, with 200 terms, proves that the root Zmin is indeed the smallest in abs value for the complete 10000th degree polynomial:

- PROOT finds all the roots of the given polynomial (here the 200th degree one), there is no missing one,
- we can be sure that the root Zmin above is also a root of the 10000th degree polynomial, in the limits of the numerical accuracy, because the terms Pn. $Z \min \wedge n$ for $n>200$ will have a modulus less than about $2 \mathrm{E}-16$ and will not contribute (again in the limits of the numerical accuracy) to the value of the complete polynomial,
- the complete polynomial can not have a smaller root Zx , because it would also be a root (in the numerical accuracy limits...) of the 200th degree polynomial, for the same reason that the $Z x^{\wedge} n$ terms for $n>200$ would be negligible.

Here is my HP-71B program using both the MATH and JPC ROMs, for reference, I didn't check but it should run in a few hours on a physical HP-71B. The HP48 and later series that have equivalent PROOT commands may find the solution in less time due to the faster CPU.

```
10 ! SRC12B2
20 OPTION BASE 0
30 !
```

```
N=200
DIM P(N)
COMPLEX Z(N)
    ! build the polynomial
    P(N)=2
    FOR I=N-1 TO 0 STEP -1
    P(I) =FPRIM (P(I+1)+1,20000)
    NEXT I
    !
    MAT Z=PROOT(P)
    A=MAXREAL
    FOR I=0 TO N-1
        A=MIN(A,ABS(Z(I)))
    NEXT I
    DISP A
    END
J-F
C.Ret 8

Posts: 223
Member Joined: Dec 2013

RE: [VA] SRC \#012b - Then and Now: Root

\section*{J-F Garnier Wrote:}
(10th November, 2022 10:26)
My understanding is that the PROOT method, with 200 terms, proves that the root Zmin is indeed the smallest in abs value for the complete 10000th degree polynomial:
- PROOT finds all the roots of the given polynomial (here the 200th degree one), there is no missing one, - we can be sure that the root Zmin above is also a root of the 10000th degree polynomial, in the limits of the numerical accuracy, because the terms Pn. \(\mathbf{Z m i n} \wedge n\) for \(\mathbf{n} \mathbf{> 2 0 0}\) will have a modulus less than about 2E-16 and will not contribute (again in the limits of the numerical accuracy) to the value of the complete polynomial,
- the complete polynomial can not have a smaller root Zx , because it would also be a root (in the numerical accuracy limits...) of the 200th degree polynomial, for the same reason that the \(\mathbf{Z x}^{\wedge} \mathbf{n}\) terms for \(\mathbf{n}>\mathbf{2 0 0}\) would be negligible.

Your explanations make sense and explain why the results no longer evolve when there are enough monômes, the truly high powers of \(x\) with \(\mid x]<1\) becoming negligible at the (limited) 12-figures precision of the machine.
This also explains why using Horner's method as I do waste a lot of computation time because the summation starts with the negligible monômes which will be overwritten by the larger values of the last ones.

Concerning the PROOT instruction, I did some tests with about twenty values for \(N\) ranging from 8 to 66 . Each time, I observe that the root of minimal absolute value (of minimal norm) is store at the first position in the result vector containing the roots.
I now short end my code by removing the search loop for the minimum value at the end of the program. I only count of the PROOT program to systematically place the smallest normed root in the first \(R(0)\) position.

It may have to do with this paragraph found on page 128, section 12 of the Math Pac Owner's Manual:
The FFOTT function is global in the sense that the user is not required to supply either an initial guess or a stopping criterion; in other words, no prior knowledge of the location of the roots is assumed. The FROOT function always attempts to begin its search (iteration) at the origin of the complex plane. An annulus in the plane known to contain the smallest magnitude root of the current (original or quotient) polynomial is constructed about the origin (using five theoretical bounds) and the initial Laguerre step is rejected if it exceeds the upper limit of this annulus. In this case, a spiral search from the lower radius of the annulus in the direction of the rejected initial step is begun until a suitable initial iterate is found.

Once the iteration process has successfully started, circles around each iterate are constructed (using two theoretical bounds) that are known to bound the root closest to that iterate; the Laguerre step size is constantly tested against the radii of these circles and modification of the step is made when it is deemed to be too large or when the polynomial value does not decrease in the direction of the step. For this reason, the roots are normally found in order of increasing magnitude, thus minimizing the roundoff errors resulting from deflation.
Having the minimal absolute value (minimal normed or smallest magnitude root) at \(R(0)\) is not just a coincidence...
EDITED: See next post where Chan show me that It is effectively just by coïncidence !!

I put a new set of AAA cells in my HP-71B to replace the previous old set that just died during the last endless attempt.

It took nearly two hours to find the minimum norm \(\boldsymbol{A B S}(\boldsymbol{R}(\mathbf{0}))=\mathbf{0 . 8 0 6 5 1 3 5 9 9 2 6 1}\) from a polynomial \(\boldsymbol{P}()\) of degree \(N=200\) with the following 3-liner directly adapted from J. -F. Garnier's code:
```

10 TO=TIME @ N=200 @ OPTION BASE O @ DIM P(N) @ COMPLEX R(N-1)
20 P(N)=2 @ P(N-1)=3 @ FOR K=N-2 TO 0 STEP -1 @ P(K)=FPRIM(2+P(K+1)) @ NEXT K
30 MAT R=PROOT(P) @ TO=TIME-TO @ DISP T0;ABS(R(0)) @ BEEP
T0
6760.46 (1:52'40.4")
ABS (R(0))
.806513599261
R(0)
(-.645758096347,-.483177676217)
ABS (R(1))
.806513599261
R(1)
(-.645758096347,.483177676217)

```
-Attached File(s)
Thumbnail(s)


\section*{Albert Chan}

Posts: 2,142
Joined: Jul 2018

\section*{RE: [VA] SRC \#012b - Then and Now: Root}

\section*{C.Ret Wrote:}
(10th November, 2022 15:29)
Concerning the PROOT instruction, I did some tests with about twenty values for \(N\) ranging from 8 to 66 . Each time, I observe that the root of minimal absolute value (of minimal norm) is store at the first position in the result vector containing the roots.

Not true. Roots obtained might not be sorted by abs.

Example, hard coded for \(N=8\), last root abs is the minimum, not first.
```

>RUN
1.81 . 735640161749
>MAT DISP R
(-.379096933597,-.630437913292) ! abs = 0.73564016175
(-.379096933597,.630437913292)
(.125003135678,-.726137429683) ! abs = 0.736818397379
(.125003135678,.726137429683)
(-.677447444085,-.29696080146) ! abs = 0.739676116352
(-.677447444085,..29696080146)
(.518497763744,.521653714169) ! abs = 0.735501548954
(.518497763744,-.521653714169)

```
C.Ret 8

Posts: 223
Member
Joined: Dec 2013

RE: [VA] SRC \#012b - Then and Now: Root
Thank you Chan.

You right, the last two or three roots may or may not have a larger norm than the first one (or first two). The loop to seek for the smallest is still needed.

Here is a corrected version of my code:
```

10 INPUT "n=";N @ TO=TIME @ OPTION BASE O @ DIM P(N) @ COMPLEX R(N-1)
20 P(N)=2 @ P(N-1)=3 @ FOR K=N-2 TO 0 STEP -1 @ P(K)=FPRIM(2+P(K+1)) @ NEXT K
30 MAT R=PROOT(P) @ FOR K=1 TO N-1 @ IF ABS(R(K))<ABS(R(0)) THEN VARSWAP R(0),R(K)
40 NEXT K @ TO=TIME-TO @ DISP TO;ABS(R(0)) @ BEEP

```

I also have to edit my previous post to indicate my mistake...

12th November, 2022, 02:59


\section*{Valentin Albillo 8}

Posts: 958
Senior Member
Joined: Feb 2015
Warning Level: 0\%

RE: [VA] SRC \#012b - Then and Now: Root
Hi, all,

Thanks for your interest in this second part of my SRC\#12 and most definitely for your solutions and comments, much appreciated as always.

I'll post my original solution next Sunday circa \(23: 00 \mathrm{GMT}+1\) (Spain is physically a GMT+0 country but for some retarded obsolete law we're stuck with \(G M T+1\), which badly shifts the times for everything,) so if anyone wants to have a further say on the subject this is the last chance.

By the way, C.Ret, you recently posted this code:

\section*{C.Ret Lately Wrote:}

INPUT " \(\mathrm{n}=\mathrm{=}: \mathrm{N}\) @ TO=TIME @ OPTION BASE 0 @ DIM \(\mathrm{P}(\mathrm{N})\) @ COMPLEX R(N-1) [...]
whis is fine save for the fact that I stated this in my \(O P\) :

\section*{... but Valentin Albillo Previously Wrote: \\ Your program should have no inputs [...]}
which disqualifies this code of yours as a valid solution. As ABBA said: "Rules must be obeyed".

Also, and it applies to everyone (me included), I think that it's proper etiquette to post not just the code per se, but also a run of it, with results and timings. This is what I always do as I feel that's the proper way to post a solution: code, run, results, timings.

Finally, I feel that few of you heeded my advice about properly balancing the work done by the program vs. the work done by the programmer, with some of you actually doing a lot of work to then have the program doing the very minimum work necessary, or the other way around.

Enough. As I said, my original solution will be posted next Sunday. Best regards.

\section*{v.}

C.Ret 8

Posts: 223
Member

\section*{RE: [VA] SRC \#012b - Then and Now: Root}

By the way, C.Ret, you recently posted this code:

\section*{C.Ret Lately Wrote:}
```

10 INPUT "n=";N @ TO=TIME @ OPTION BASE O @ DIM P(N) @ COMPLEX R(N-1) [...]

```
whis is fine save for the fact that I stated this in my \(O P\) :

\section*{... but Valentin Albillo Previously Wrote:}

Your program should have no inputs [...]
which disqualifies this code of yours as a valid solution. As ABBA said: "Rules must be obeyed"

That's life, every mistake has to be paid cash!

I copy-paste the wrong version, ...
... Now I am disqualified (but was I able to solve this quizz - not sure really )
... Now I am in big trouble with my internal Q\&C, the voll team at the Research Département and HQ directors since I post at a public place a true confidential documentation and badly secured it.
despite I am out of the race now, I still looking for a better solution...

Still have a question, I notice that all the roots found by PROOT are within the unit circle centered at the origin of the complex plane.
Is there any reason for that?
Is there a relation between this all-prime coefficient polynôme and something with a circle or a trigonometry fact? Isn't there a quick and efficient link between its roots and any characteristic trigonometric fUnction?
hummm.

\section*{J-F Garnier 8}

Senior Member

Posts: 790
Joined: Dec 2013

RE: [VA] SRC \#012b - Then and Now: Root

\section*{C. Ret Wrote:}
(12th November, 2022 07:21)
Still have a question, I notice that all the roots found by PROOT are within the unit circle centered at the origin of the complex plane.
Is there any reason for that?

I believe for the same reason I indicated in my first post above:

\section*{J-F Garnier Wrote:}
(9th November, 2022 10:43)
My first thoughts: a real root must be negative, and greater than -1 since the polynomial quickly takes very large values for \(|X|>1\).

Transposed for complex roots, it means the roots must be \(|z|<1\).
And I believe this is also related to this comment of Valentin:

\section*{Valentin Albillo Wrote:}
(12th November, 2022 02:59)
Finally, I feel that few of you heeded my advice about properly balancing the work done by the program vs. the work done by the programmer, with some of you actually doing a lot of work to then have the program doing the very minimum work necessary, or the other way around.

This is experimental math for most of us, as Valentin recently pointed it out. Fernando and I (and maybe others silently) tried with PROOT and found a candidate for the minimum root, and the fact that it was small enough lead me to my analysis and my solution with 200 terms.

J-F

\section*{Albert Chan}

Senior Member

Posts: 2,142
Joined: Jul 2018

\section*{RE: [VA] SRC \#012b - Then and Now: Root}

For \(P(x)\) degree 24 or higher, min abs root is around \(z=\left(-2 \pm \sqrt{ } 2 *_{i}\right) / 3\)
We can use Newton's method to zeroed in true root, for \(N \geq 24\)
\(|z|^{\wedge} 2 \approx(4+2) / 9=2 / 3-->|z|^{\wedge} 148 \approx(2 / 3)^{\wedge} 74 \approx 9 E-14\)

Set \(N=148\), we have 12-digits accuracy for \(N \geq 148\) min abs root.

10 DESTROY ALL @ SETTIME 0
\(20 \mathrm{~N}=148\) @ DIM \(\mathrm{P}(\mathrm{N})\) @ \(\mathrm{P}(1)=3\) ! odd primes
30 FOR \(K=2\) TO N @ P(K)=FPRIM(P(K-1)+2) @ NEXT K
40 COMPLEX Z,F0,F1 @ \(Z=(-2, S Q R(2)) / 3 @ A 1=-1\)
\(50 \mathrm{~A} 0=\mathrm{A} 1\) @ \(\mathrm{A} 1=\mathrm{ABS}(Z)\) @ DISP TIME,Z,A1 @ IF A1=A1+10*(A1-A0)^2 THEN END
60 F0=0 @ F1=0 @ FOR K=N TO 1 STEP -1 @ F0=F0*Z+P(K) @ F1=F1*Z+K*P(K) @ NEXT K
70 F0=F0*Z+2 @ Z=Z-F0/F1! newton's method
80 GOTO 50
\(>\) RUN
\(.51(-.666666666667, .47140452079) .816496580927\)
\(1.9(-.645842585444, .480858085479) .805193854636\)
\(3.22(-.645737371307, .48318762432) .806502965276\)
\(4.53(-.645758096686, .483177673982) \quad .806513598193\)
\(5.9(-.645758096347, .483177676218) \quad .806513599261\)

14th November, 2022, 01:23 (This post was last modified: 14th November, \(202201: 57\) by PeterP.)
Post: \#19


\section*{PeterP 8}

Posts: 172
Member
Joined: Jul 2015
RE: [VA] SRC \#012b - Then and Now: Root
This is in homage to a lot of learning that I was allowed to do here about polynomial roots, thanks to the ever generous author and Albert Chan.

I tried Barstow's method as that was something I found on the forum but the hp41 only gets to about 20-30 terms. Not enough.

However, given the 10 digit accuracy, the first 100-128 terms should be enough. Thanks to the teachings from Albert, I implemented a root squaring algorithm, looking for the max abs root of \(1 / P(x)\).

I can get it to run with about 100 terms, given the memory limitations of the HP41CX.

Result is 0.806427842 in about 23 seconds. Not very impressive accuracy. More squaring would be required, but I dont have enough registers, and not enough digits of accuracy either I guess.

The code first creates the list of 100 or so primes and stores it into registers. And then successively squares them until we get only 3 elements. And then calculates the max abs root by dividing the third element by the first element, taking the ( 2 * number_of_squaring) th root, for the max abs of \(Q(x)=1 / P(x) .1 / x\) gives then the min abs root.

The listings are attached as pictures. I will try to type them up as well, but I am worried about typos.

Thank you again for a wonderful learning experience!
Cheers
Peter

Attached File(s)
Thumbnail(s)


14th November, 2022, 01:44 (This post was last modified: 14th November, 2022 01:55 by PeterP.)


Posts: 172
Member
Joined: Jul 2015

RE: [VA] SRC \#012b - Then and Now: Root

Ok, here are the listings typed up, please excuse any typos:

It is called with
XEQ "VA2" and stops after about 23 seconds with the result in X .
The elements \(a[n]\) to \(a[n-x]\) are in the registers sss.eee, with sss \(>\) eee. This control number is stored in R00
The calculated elements \(b[n]\) from the squaring are in registers \(b b b\).fff with \(b b b>f f f\). This control number is stored in R01
Once all \(b[n]\) are calculated, the new boundaries are calculated (1/2 of bbb-fff. And bbb-fff is half of sss-eee) The \(b[n]\) 's become the new \(a[n] s\), and the control numbers are created accordingly,

R02 accumulates the values for each \(b[n-i]\) as the summation over all \(j=1 \ldots . i\) takes place.
R03 holds the current \(j\)
R04 holds the current i
R06 holds the number of elements to calculate
LBL 00 is the loop summing over the elements \(j=1 \ldots\) i
LBL 01 is the loop over all \(i\)
LBL 03 is the loop over the successive \(b[n], b^{\prime}[n], b^{\prime \prime}[n]\), etc
LBL 05 is the final min abs root value calculation after the last \(\mathrm{b}^{\prime \prime} \ldots\) [n] series calculation which has only 3 elements.

REM - Fill Primes in descending order into registers. Called with a control word of sss.eee

LBL "PL"
50
SETCSPD
RDN
STO M
2
STO IND M
DSE M
VIEW X
3
VIEW X
STO IND M
DSE M
LBL 00
INCX
INCX
PRIME?
GTO 01
LASTx
GTO 00
LBL 01
View X
STO IND M
DSE M
GTO 00
END
REM..... Main PRogram
LBL "VA2"
109.009

STO 00
XEQ "PL"
RCL 00
159.110

STO 01
X<>Y

INT
STO 06
LBL 03
1
STO 04
RCL Ind 00
X^2

STO IND 01
LBL 01
RCL 04
STO 03
CLX
STO 02
LBL 00
RCL 03
RCL 04
RCL 00
\(+\)
RCL IND \(X\)
STO 05
RCL 03
CHS
RCL 04
RCL 00
\(+\)
RCL IND X
ST* 05
2
-1
RCL 04
RCL 03
\(Y^{\wedge} x\)
*
RCL 05
*
\(\mathrm{ST}+02\)
DSE 03
GTO 00
RCL 00
RCL 04
RCL IND X
\(X^{\wedge} 2\)
-1
RCL 04
Y^X
*
RCL 02
\(+\)
RCL 01
RCL 04
\(X<>Y\)
STO IND Y
RCL 04
INCX
STO 04
Enter
RCL 06
\(X>Y\) ?
GTO 01
RCL 06
2
/
INT
STO 06
2
\(X>Y\) ?
GTO 05
RDN
RCL 00
INT
\(X<>Y\)
1000

RCL 00
INT
\(+\)
X<>01
X<>00
GTO 03
LBL 05
-2
RCL 01
RCL IND X
RDN
\(+\)
RCL IND X
\(\mathrm{R}^{\wedge}\)
/
RCL IND 01
LOG
2
LOG
/
2
*
1/x
\(y^{\wedge} X\)
1/x
Beep
CLD
END

\section*{Albert Chan 8}

Posts: 2,142
Senior Member

\section*{RE: [VA] SRC \#012b - Then and Now: Root}

It may help to explain PeterP's root squaring program.
This was a PM I sent to PeterP, and other members.

\section*{Albert Chan Wrote:}

Instead of min abs \(P\) roots, we solve for max abs \(Q\) root, \(Q(x)=P(1 / x)\)
Graeffe's root squaring method (next row roots \(=\) previous roots^\({ }^{\wedge} 2\) )
We only show top 3 coefficients, because we only care for Q max abs root
\(\mathrm{q}=[2,3,5,7,11,13,17,19,23,29, \ldots] / /\) assumed infinite degree polynomial
\(\rightarrow\left[2^{\wedge} 2,11,27\right]\)
\(\rightarrow[2 \wedge 4,95,303]\)
\(\rightarrow[2 \wedge 8,671,7775]\)
\(\rightarrow\left[2^{\wedge} 16,3.530559 \mathrm{E} 6,1.01291839 \mathrm{E} 8\right]\)
\(\rightarrow\) [2^32, 8.11677068927E11, 4.099840909585279E15]
\(\rightarrow[2 \wedge 64,3.455854558672141 \mathrm{E} 25,1.816641529875401 \mathrm{E} 31]\)
\(\rightarrow\left[2^{\wedge} 128,-5.240706455638273 E 50,2.748844184968018 \mathrm{E} 62\right]\)
\(\rightarrow\left[2^{\wedge} 256,8.757340043013178 \mathrm{E} 100,7.559454552508339 \mathrm{E} 124\right]\)
\(\rightarrow\) [2^512, 9.837400259693430E201, 5.714553957282261E249]
...
2nd column are not "pure squares" (note the negative sign), but 3rd column is.
Thus, roots to seek are complex conjugates.
Assuming roots are well separated now:
\(\max\) abs Q root \(=\operatorname{surd}\left(5.714553957282261 \mathrm{E} 249 / 2^{\wedge} 512,512 * 2\right)=1.2399046971021601\)
\(\min\) abs \(P\) root \(=1 /(\max\) abs \(Q\) root \()=0.8065135992606103\)
[VA012b] is similar to [VA012a], a diffusion problem.
Root Squaring process also based from its neighbor cells.
\(\mathrm{b}[\mathrm{n}]=\mathrm{a}[\mathrm{n}] \wedge 2\)
\(b[n-1]=-a[n-1]^{\wedge} 2+2 * a[n] * a[n-2]\)
\(b[n-2]=+a[n-2]^{\wedge} 2-2 * a[n-1]^{*} a[n-3]+2 * a[n] * a[n-4]\)

By the time big primes effect diffused to the top, its effects are miniscule

\section*{\(\rightarrow\) EMAIL PM \(\sim\) FIND}

QUOTE की REPORT

\section*{Valentin Albillo 8}

Posts: 958
Joined: Feb 2015
Warning Level: 0\%

RE: [VA] SRC \#012b - Then and Now: Root

\section*{Hi, all,}

Well, a full week has elapsed since I posted my \(O P\), which has already passed the 1,300 views mark (and Problem 1 has exceeded 9,000 views already,) and I've got a number of solutions and/or comments, namely by Werner, Jean-
François Garnier, C.Ret, PeterP, Fernando del Rey and Albert Chan. Thank you very much to all of you for your interest and valuable contributions.

Now I'll provide my original solution to this Problem 2 but first a couple' comments:
1) Some of you provided just code but no numerical results, others provided numerical results but no code, and most of you provided comments but didn't provide timings. As I said in my previous post, I'd consider proper etiquette and most useful for everyone if contributors would kindly post program code, a sample run, results and timings. Not a rule but it would help. And comments are always most welcome, of course.
2) Again, I'm mildly surprised that no one posted RPL solutions or RPL code of any kind. As for RPN code, PeterP made a most brave attempt to get a reasonably accurate result (as did Werner with Free42, though using capabilities not available on the vintage HP-42S, ) but it might be that RPL people deem this challenge as too trivial for their powerful vintage RPL calcs. On the other hand, perhaps it might be that ... naw, never mind.

That said, these are my own approach and resulting original solution:

First of all, the main difficulty is the \(\mathbf{1 0 , 0 0 0}{ }^{\text {th }}\)-degree. Were it a mere \(100^{\text {th }}\)-degree polynomial, it would be quite trivial, but dealing with the full polynomial using PROOT, say, would require allocating 10,000 real elements for the coefficients (they don't fit as integers) plus another 10,000 complex elements for the roots, for a grand total of \(\sim 30,000 \times 8=\mathbf{2 4 0}\) \(\mathbf{K b}\), to which you must add the considerable memory that PROOT needs internally ( \(21 \times 10.000+261 \sim 210 \mathbf{K b}\), ) which adds up to \(\sim \mathbf{4 5 0} \mathbf{~ K b}\), surely exceeding maximum available RAM).

Then again, finding the 10,000 roots woud take \(\sim(10,000 / 100)^{2} x\) the time required to deal with a \(100^{\text {th }}\)-degree polynomial, which on a physical HP-71B is about 2,100 sec., so the big one would take \(2,100 \times 100^{2} \sim 243\) days. In other words, utterly unfeasible. So much for sheer brute force ...

Now, from theoretical considerations it's pretty obvious that none of the 10,000 roots can have an absolute value (aka modulus, aka magnitude) \(>\mathbf{1}\), because then the highest term, \(\mathbf{1 0 4 , 7 4 3} \mathbf{x} \mathbf{1 0 , 0 0 0}\), would dominate the sum, making it nonzero. Likewise, if the roots are too close to \(\boldsymbol{O}\) then their powers will quickly tend to zero, thus failing to ever contribute enough negative value to compensate for the lowest coefficient, 2. Thus, all the roots must reside in an annulus of external radius \(\mathbf{1}\) and internal radius to be determined by the minimum magnitude among all the roots, which we can roughly estimate as I'll explain in a moment.

Once we have a rough estimation for said minimum magnitude, we can compute another estimation, this time for the minimum degree of the truncated polynomial so that its highest term already contributes negligibly to the evaluation of the polynomial according to the precision required, say 12 digits.

My resulting program is thus this 5-liner: (REPEAT/UNTIL are from the JPC ROM, just for show)
```

1 DESTROY ALL @ @ OPTION BASE O @ M=0 @ W=O @ N=0 @ FIX 2
2 REPEAT @ W=M @ N=N+10 @ GOSUB 4 @ DISP N;M;P @ UNTIL ABS (M-W)<.01 @ STD
3 N=IP((-12-LGT(P))/LGT(M)) @ DISP "Deg:";N;".." @ GOSUB 4 @ DISP "Min:";M;P @ END
DIM A(N) @ COMPLEX R(N) @ P=2 @ A (N)=P @ FOR I=1 TO N @ P=FPRIM(P+1) @ A (N-I)=P
5 NEXT I @ MAT R=PROOT (A) @ M=1 @ FOR I=0 TO N-1 @ M=MIN(M,ABS(R(I))) @ NEXT I @ RETURN

```

Line 1 does some initialization.

Line 2 finds the minimum absolute value for polynomials of degrees \(10,20,30, \ldots\) until two consecutive minimum values are within 0.01 , which essentially gives us the result correct to 2 digits ( \(\sim 0.81\) )

Line 3 now uses this 2-digit value to compute an estimation to the minimum necessary degree to get a fully accurate result, which it then obtains and displays, correct to 12 digits. It is important to use the minimum-degree truncated polynomial for speed reasons ("the faster, the better") because PROOT finds the roots of a \(153^{\text {th }}\)-degree polynomial about \(\mathbf{7 1 \%}\) faster than for a \(200^{\text {th }}\)-degree one, say.

Lines 4 and 5 are a subroutine which creates the \(N^{t h}\)-degree polynomial, fills it up with the prime coefficients, computes all \(N\) complex roots and returns the minimum magnitude among them. All of it in just 2 lines of code.

This way, the user doesn't have to guess and provide any particular degree for the truncated polynomial at all, the program finds the necessary degree and then the sought-for minimum absolute value to maximum accuracy (12 digits). In other words, the program does all the work, fast.

Let's do a sample run:
>RUN
\begin{tabular}{llll} 
Deg & Min & Ncoef & \\
---------------- & \\
10 & 0.73 & 31 & \\
20 & 0.80 & 73 & \\
30 & 0.81 & 127 & \\
40 & 0.81 & 179 & \\
Deg: & 153 & & \\
Min: & .806513599261 & 887
\end{tabular}
correct to 12 digits in \(\mathbf{1} \mathbf{~ h r} 18\) min on a physical HP-71B. We can check the highest term contribution like this:
```

>P;M;N;P*M^N -> 887 . 806513599261 153 4.56575832163E-12

```
which indeed contributes negligibly to the value of the polynomial so it and all other higher terms can be safely ignored.

As for using РROOT to find all the \(N\) roots (most of them complex) vs. using Newton's method to directly find the root that results in the desired minimum magnitude (as A.Chan does in post\#18,) the reason is that in general that might not guarantee that you get the absolute minimum.

In A.Chan's post it works but in general it could be just luck, no guarantee at all without further considerations to try and find a working initial guess, and as you can see in the graph below most roots are within a thin annulus of radii \(\mathbf{1}\) and \(\mathbf{0 . 8 2}\), i.e. only 0.18 wide, and the conjugated pair we need are extremely near to it, so the roots are quite clustered and isolating the one we want might be a hit-or-miss affair, while using PROOT and then finding the minimum magnitude is \(100 \%\) guaranteed to find the correct one, giving us perfect ease of mind with minimum effort on our part.


Well, that will be it for now. Thanks again to all of you who contributed and some of you even solved both Problem 1 and Problem 2, a perfect 2 for 2 score so far, but as they were the easiest problems of the lot I wonder if you'll manage to also solve incoming Problem 3, which is sure to raise the bar ... slightly ? We'll see ...

\section*{v.}

\section*{Albert Chan}

Posts: 2,142
Senior Member
RE: [VA] SRC \#012b - Then and Now: Root
If we consider polynomial coefficients with geometric progression:
\(f(x)=1+\left(r^{*} x\right)+\left(r^{*} x\right)^{\wedge} 2+\ldots+\left(r^{*} x\right)^{\wedge} n=\left(\left(r^{*} x\right)^{\wedge}(n+1)-1\right) /\left(\left(r^{*} x\right)-1\right)\)
From RHS numerator, all f roots \(a b s=1 / r\)
\(P(x)=2+3 x+5 x^{\wedge} 2+7 x^{\wedge} 3+11 x^{\wedge} 4+\ldots\)

P roots, sorted in abs: ( \(-0.6458 \pm 0.4832 i),(0.4472 \pm 0.7248 i),(-0.2853 \pm 0.8292 i), \ldots\) With corresponding abs: 0.8065, \(0.8517,0.8769, \ldots\)

Primes does not grow as fast as geometric progression. (sum of reciprocal primes diverges)
\(P\) min abs root is due to ratio, \((5 / 2=2.5)>(11 / 5=2.2)\)
If the ratios were about the same, we expected \(f(x)\) roots pattern, with similar sized roots.

Example, \(R(x)=P(x)+0.5,5 /(2+0.5)=2\).
\(R\) roots, sorted in abs: ( \(-0.6811 \pm 0.5122 i),(0.4692 \pm 0.7114 i),(-0.2905 \pm 0.8666 i), \ldots\)
With corresponding abs: 0.8522, 0.8522, \(0.9140, \ldots\)

If guess close enough, we can use Newton's method to zeroed in P min abs root.
```

EMAIL PPM O FIND

```

\section*{Valentin Albillo 8}

Senior Member

RE: [VA] SRC \#012b - Then and Now: Root

\section*{Hi, Albert Chan,}

Thanks for your recent additional comments, I appreciate it. However, I have a thing or two to comment back, read on ... (all highlights are mine)

\section*{Albert Chan Wrote:}

\section*{Valentin Albillo Wrote:}

Line 2 finds the minimum absolute value for polynomials of degrees 10, 20, 30, \(\ldots\) until two consecutive minimum values are within 0.01 , which essentially gives us the result correct to 2 digits ( \(\sim 0.81\) )

\section*{Slight error in logic.}

Not at all, see below.

\section*{Albert Chan Wrote:}

If consecutive minimum abs both \(\sim 0.81\), we cannot deduce trend apply to higher degree. (we can assume trend continues, but have to later test validity of assumption)

And I did test, it's just that I didn't want to make an already long post any longer by including unneeded data that most people won't be interested in, as they understand the scope of my articles and challenges and trust my results (which they can verify by themselves, if in doubt), but as you seem to like said data, here you are, the checks I did (which I saved but didn't post):
```

Degree Min. Magnitude
-----------------------
1 . .66666 6666667
2 . 63245 5532034
4 . 65224 6975033

```
\begin{tabular}{rlll}
8 & .73550 & 1548954 \\
16 & .76890 & 4440166 \\
32 & .80252 & 6072477 \\
64 & .80650 & 0035750 \\
96 & .80651 & 362173 \\
128 & .80651 & 3599285 \\
& & \\
140 & .80651 & 3599258 \\
145 & .80651 & 3599260 \\
150 & .80651 & 3599261 \\
160 & .80651 & 3599261 \\
165 & .80651 & 3599261 \\
170 & .80651 & 3599261 \\
175 & .80651 & 3599261 \\
180 & .80651 & 3599261 \\
185 & .80651 & 3599261 \\
190 & .80651 & 3599261 \\
195 & .80651 & 3599261 \\
200 & .80651 & 3599261
\end{tabular}
and I was more than satisfied that the trend continued alright and converged to the correct solution, namely . 806513599261.

\section*{Albert Chan Wrote:}

\section*{Valentin Albillo Wrote:}

Line 3 now uses this 2-digit value to compute an estimation to the minimum necessary degree to get a fully accurate result, which it then obtains and displays, correct to 12 digits.

Close, but not quite.
If this "minimum necessary degree" polynomial also gives abs \(\leq 0.81\), we are done.
However, if it gives bigger abs, we have to repeat again, with even higher degree polynomial.

But it didn't give "bigger abs", it gave . 806513599261, which is less than 0.81 , so your comment doesn't apply at all and my statement is fully correct, not merely "close".

Now a word on the scope for my articles and challenges. In a past thread in which you took part (you posted 5 times no less) and so you surely read my posts there, I said:
"My articles are intended as just that, articles to be published in a physical fan-made magazine [...] or else on the WWW (MoHPC) for all kinds of fans, most of them not scholars, so my articles do not have the structure nor goals of a formal peer-reviewed paper."

In other words, my goal is first and foremost to provide entertainment to the forum members and HP calc fans in general, enticing them to think about the challenge and how to use their vintage HP calc to solve it, and if additionally they learn something new and interesting (and even useful) from my productions then so much the better. That is my goal.

Yours is obviously different, seemingly centered on lecturing, posting symbolic proofs and theoretical ramblings and lots of data obtained in Xcas sessions, lua, Mathematica, the works. Good for you and for the people who like (lots of) posts like that. Not my cup of tea here.

And please leave aside topics having little or nothing to do with my present Problem 2 (the references to integration, Borwein integrals and Kahan), save that for your own threads or where it's appropriate. Thanks.

\section*{v.}


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